### Learning in Big Data Analytics Lecture 6

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### Direct Discovery of Communities



### INTRODUCTION

- ► So far, we partitioned graphs into disjoint communities
- But communities might be overlapping
- Solution: Determine communities as (induced) subgraphs of a certain type
- Subgraphs should contain unusually large amount of edges
- ► Will treat two types briefly here:
  - Cliques
  - Complete bipartite subgraphs



### FINDING CLIQUES

DEFINITION [INDUCED SUBGRAPH] Let G = (V, E) be a graph. A subgraph  $C = (V' \subset V, E' \subset E)$  is *induced* iff  $(v', w') \in E$  implies  $(v', w') \in E'$ 

for any  $v', w' \in V'$ .

DEFINITION [CLIQUE]

Let G = (V, E) be a graph.

- An induced subgraph C = (V', E') is called a *clique* iff any pair of nodes in *C* is connected by an edge.
- ► A clique C = (V', E') is *maximal* iff extending the clique by any node and its edges implies that the clique property no longer holds.



### Communities as Cliques

- Possible idea: Determine communities as maximal cliques
- *Caveat:* The number of maximal cliques in a graph may be exponential in the number of nodes
- So, listing all maximal cliques is a computationally demanding problem
- Nevertheless, identifying communities as clique like arrangements is popular



### COMPLETE BIPARTITE GRAPHS

DEFINITION [(COMPLETE) BIPARTITE GRAPHS]

A graph G = (V, E) with vertices V and edges E is referred to as *bipartite* iff

• there are  $V_1, V_2 \subset V$  such that

 $V = V_1 \cup V_2$  and  $E \subset (V_1 \times V_2)$ 

• A bipartite graph G = (V, E) is *complete* iff

 $V = V_1 \stackrel{.}{\cup} V_2$  and  $E = (V_1 \times V_2)$ 

that is iff each node from  $V_1$  is connected with each node from  $V_2$ 

- A complete bipartite graph where  $|V_1| = s$ ,  $|V_2| = t$  is referred to as  $K_{s,t}$
- A complete bipartite graph is also referred to as *biclique*



### COMPLETE BIPARTITE GRAPHS AND COMMUNITIES

- ► *Strategy:* Seek to discover all sufficiently large bicliques
- ► Treat them as "nuclei" (or seeds) of communities
- *Theoretical Advantage over Cliques:* While it is not possible to guarantee the existence of large cliques for graphs with many edges, one can guarantee the existence of large bicliques



### FINDING COMPLETE BIPARTITE GRAPHS

Frequent Itemset Mining Problem

- ► Let G = (V, E) on  $V = V_1 \cup V_2$  be a (large) bipartite graph
- Items are nodes from  $V_1$
- ► Baskets are nodes from *V*<sub>2</sub>
- ▶ Items in baskets are nodes from *V*<sup>1</sup> connected to basket node
- $K_{s,t}$  in *G* is itemset of size *s* that appears in *t* baskets
- So mining for frequent itemsets at threshold *t* dicovers all  $K_{s,t}$



### The Graph Affiliation Model



### **OVERLAPPING COMMUNITIES**



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- *Observation:* Communities in social networks can overlap
- Graph partitioning does not help in these cases

► Would like to have a statistical interpretation of network data

## NONOVERLAPPING VERSUS OVERLAPPING COMMUNITIES



Left: Nonoverlapping communities Right: Overlapping communities Adopted from mmds.org

- Communities may overlap or not
- Issue: How to determine communities correctly?





#### Networks and their adjacency matrices

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- ► Left: No overlap, adjacency matrix sparse across communities
- Middle: Loose overlap, adjacency matrix less sparse in shared part
- Right: Tight overlap, adjacency matrix dense in shared part



### COMMUNITY DISCOVERY: GOAL



Revealing (overlapping) communities

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- ► *Goal:* Discover communities correctly
- Regardless of whether they overlap or not

Determine the statistically most likely community structure



- ► *Issue:* Statistical control over community structure of a network
- ► Idea: Design generative probability distribution
- Given a number of nodes, this generative distribution generates edges
- The generative distribution represents a particular community structure
  - The distribution knows about nodes belonging to communities
  - It generates more edges within communities
  - It generates less edges between communities



The generative distribution represents community structures

- The distribution knows about nodes belonging to communities
- It generates more edges within communities
- It generates less edges between communities



Distribution representing a community structure generating network

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Distribution representing a community structure (left) generating network (right) Adopted from mmds.org

- ► We can generate networks when knowing community structure
- ► *But:* We would like to determine the community structure when knowing the network

Isn't that exactly the opposite?



### GENERATIVE DISTRIBUTIONS



We can do this: generating network from distribution...

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...but we want this: inferring distribution from network

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# GENERATIVE DISTRIBUTIONS: MAXIMUM LIKELIHOOD INFERENCE



We want to infer distribution from network

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Maximum Likelihood Estimation

- Let Θ be a *parameterized class of probability distributions* that generate networks
  - We identify the different distributions with the different parameterizations
     Formally not 100% correct, but doesn't matter here
- ► Let  $\mathbf{P}(N \mid \theta)$  be the probability that distribution  $\theta \in \Theta$  generates network *N*



# GENERATIVE DISTRIBUTIONS: MAXIMUM LIKELIHOOD INFERENCE



We want to infer distribution from network

Adopted from mmds.org

Maximum Likelihood Estimation

- ► Let  $\mathbf{P}(N \mid \theta)$  be the probability that distribution  $\theta \in \Theta$  generates network *N*
- Maximum likelihood estimation: Determine distribution θ̂ that generated N with greatest likelihood:

$$\hat{\theta} := \operatorname*{arg\,max}_{\theta \in \Theta} \mathbf{P}(N \mid \theta) \tag{1}$$

UNIVERSITÄT his computes most reasonable distribution  $\hat{\theta}$  for network N belefeld

### AFFILIATION GRAPH MODEL: DEFINITION I

- An AGM θ generates a network N = (V, E) by adding edges E to a given set of nodes V
- ► For  $u, v \in V$ , edge (u, v) is generated with probability  $\mathbf{P}_{\theta}((u, v))$
- $\mathbf{P}_{\theta}((u, v))$  depends on the parameters  $\theta$
- Recall that  $\theta$  specifies community structure

So, what exactly is  $\theta$  supposed to be?



### AFFILIATION GRAPH MODEL: PARAMETERS

- C, as a set of *communities*
- $M \in \{0,1\}^{C \times V}$ , specifying assignment of nodes  $v \in V$  to communities  $C \in C$ , where

$$M_{C,v} = \begin{cases} 1 & v \text{ belongs to } C \\ 0 & \text{otherwise} \end{cases}$$
(2)

- *M* specifies "affiliations" of nodes  $v \in V$
- Note that one can vary C, as a parameter, but not V
- ►  $(p_C)_{C \in C}$  as probabilities to generate edges (u, v) because  $u, v \in C$
- Summary: A particular AGM  $\theta$  corresponds to

$$\theta = (\mathcal{C}, M, (p_C)_{C \in \mathcal{C}}) \tag{3}$$



**Several** *C* **containing both** *u*, *v* 

- Let  $M_u, M_v \subset C$  be the subsets of communities that contain u and v, respectively
- Existence of communities that contain both *u*, *v* means

 $M_u \cap M_v \neq \emptyset$ 

- Memberships in different communities have no influence on each other
- ► That is, we assume *statistical independence*



Several C containing both u, v

Statistical independence is expressed by

$$\prod_{C \in M_u \cap M_v} (1 - p_C)$$

as probability of *no edge* (u, v) *in any community*  $C \in M_u \cap M_v$ 

• Hence, the probability to generate (u, v) is

$$1 - \prod_{C \in M_u \cap M_v} (1 - p_C) \tag{4}$$

**Done? No:** What about 
$$M_u \cap M_v = \emptyset$$
?



#### **No** *C* **containing both** *u*, *v*

For  $M_u \cap M_v = \emptyset$ , computing (4) yields (empty product is 1)

$$1 - \prod_{C \in \emptyset} (1 - p_C) = 1 - 1 = 0$$

- No edges across communities makes no sense
- Let  $\epsilon > 0$  be small; we generate an edge (u, v) with probability

$$\mathbf{P}_{\theta}((u,v)) = \epsilon \quad \text{if} \quad M_u \cap M_v = \emptyset$$



AFFILIATION GRAPH MODEL (AGM)

• An edge (u, v) is generated with probability

$$\mathbf{P}_{\theta}((u,v)) = \begin{cases} 1 - \prod_{C \in M_u \cap M_v} (1 - p_C) & M_u \cap M_v \neq \emptyset \\ \epsilon & M_u \cap M_v = \emptyset \end{cases}$$
(5)

- Edges (u, v) are generated independently from one another
- *Overall:* The probability  $\mathbf{P}_{\theta}(E)$  to generate edges *E* given AGM  $\theta$  computes as

$$\mathbf{P}_{\theta}(E) = \prod_{(u,v)\in E} \mathbf{P}_{\theta}((u,v)) \times \prod_{(u,v)\notin E} 1 - \mathbf{P}_{\theta}((u,v))$$
(6)

where  $\mathbf{P}_{\theta}((u, v))$  are computed following (5), with  $\theta = (\mathcal{C}, M, p_{\mathcal{C}})$  determining  $p_{\mathcal{C}}$  and  $M_u, M_v$  and so on.



### AFFILIATION GRAPH MODEL: OVERALL PROBABILITY

AFFILIATION GRAPH MODEL (AGM)

• The probability  $\mathbf{P}_{\theta}(E)$  to generate *E* given  $\theta$  is

$$\mathbf{P}_{\theta}(E) = \prod_{(u,v)\in E} \mathbf{P}_{\theta}((u,v)) \times \prod_{(u,v)\notin E} 1 - \mathbf{P}_{\theta}((u,v))$$
(7)

• *Reminder:* For a given network N = (V, E), the *goal* is to determine

 $\hat{\theta} := \operatorname*{arg\,max}_{\theta \in \Theta} \mathbf{P}_{\theta}(E)$ 

• That is, we need to vary  $\theta = (C, M, p_C)$  until  $\mathbf{P}_{\theta}(E)$  is maximal

How to systematically vary  $\theta = (C, M, p_C)$ ?



ISSUES

- Search space of combinations of
  - ► Communities *C*,
  - ► Assignments of nodes to communities *M*, and
  - Probabilities *p*<sub>C</sub> for communities

tends to be huge

- Concise formulas of (7) for  $\mathbf{P}_{\theta}(E)$  as function of  $\theta$  too difficult
- ► Analytical solution for determining \(\heta\) := arg max<sub>\(\theta\) \in \OPE\)</sub> P<sub>\(\theta\)</sub>(E) not available
- Moreover, parameters are both discrete (C, M) and continuous (( $p_C$ )<sub> $C \in C$ </sub>)



Approach

- 1. Pick initial set of parameters  $\theta_0$
- 2. Vary  $\theta$  such that  $\mathbf{P}_{\theta}(E)$  iteratively increases
- 3. Vary C or M first

Partial derivates of  $\mathbf{P}_{\theta}(E)$  wrt.  $p_{C}$  computable on fixed C, M

- 4. Determine optimal  $(p_C)_{C \in C}$ , e.g. by gradient descent
- 5. Keep change if  $\mathbf{P}_{\theta}(E)$  has increased, discard otherwise



Iterative variations of  $\mathcal{C}, M$ 

- ► Varying M:
  - Delete node from community, i.e. for  $M_{C,v} = 1$ , set  $M_{C,v} = 0$
  - Add node to community, i.e. for  $M_{C,v} = 0$ , set  $M_{C,v} = 1$
- ► Varying C:
  - Merge two communities
  - Split community
  - Delete community
  - Add new community, with initial random selection of members



SOFT COMMUNITY MEMBERSHIP

- ▶ Instead of  $M_{C,v} \in \{0,1\}$ , allow any real-numbered  $M_{C,v} \ge 0$
- For (u, v) to be generated because of  $u, v \in C$ , let

$$\mathbf{P}_{\theta}((u,v)) = 1 - e^{-M_{C,u}M_{C,v}}$$
(8)

be the individual probability

Proceeding exactly as before, we obtain

$$\mathbf{P}_{\theta}(E) = \prod_{(u,v)\in E} (1 - e^{-\sum_{C} M_{C,u} M_{C,v}}) \prod_{(u,v)\notin E} e^{-\sum_{C} M_{C,u} M_{C,v}}$$
(9)



SOFT COMMUNITY MEMBERSHIP

► Probability for edges *E*:

$$\mathbf{P}_{\theta}(E) = \prod_{(u,v)\in E} (1 - e^{-\sum_{C} M_{C,u} M_{C,v}}) \prod_{(u,v)\notin E} e^{-\sum_{C} M_{C,u} M_{C,v}}$$
(10)

- On fixed communities, include *M* in gradient descent (or related) optimization step
- ► Advantages:
  - Only one gradient descent run necessary
  - Less prone to get stuck in unfavorable local optima
- ► If necessary, add or delete communities, and re-run



### GENERAL / FURTHER READING

### Literature

Mining Massive Datasets, Sections 10.3, 10.5 http://infolab.stanford.edu/~ullman/mmds/ ch10.pdf

